

# Variants of Average Consensus

Yilin Mo

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## 1 Finite Time Average Consensus

Consider the following update equation:

$$x(k+1) = (I - \alpha(k)L)x(k), \quad (1)$$

We know if we fix  $0 < \alpha(k) < 2/\lambda_n(L)$ , then

$$\prod_{k=0}^{\infty} (I - \alpha(k)L) = J = \frac{\mathbf{1}\mathbf{1}^T}{n}.$$

Assume that  $L = U\Lambda U^T$ , then

$$\prod_{k=0}^{n-2} = U \begin{bmatrix} 1 & & & \\ & f(\lambda_2(L)) & & \\ & & \ddots & \\ & & & f(\lambda_n(L)) \end{bmatrix} U^T,$$

where  $f(x)$  is an  $n-1$ th degree polynomial of the following form:

$$f(x) = \prod_{k=0}^{n-2} (1 - \alpha(k)x).$$

Hence, if we choose  $\alpha(0) = 1/\lambda_2(L), \dots, \alpha(n-2) = 1/\lambda_n(L)$ , then  $f(\lambda_i(L)) = 0$ , for any  $i = 2, \dots, n$ . Thus, we can reach consensus in  $n-1$  steps.

In general, if we do not know all the eigenvalues of  $L$ , but suppose that we know  $\lambda_i(L) \in [a, b]$ , for all  $i = 2, \dots, n$ . Further assume that we can use a periodic  $\alpha(k)$ , with  $\alpha(k+T) = \alpha(k)$ , then the problem becomes finding a  $T$ th polynomial  $f(x)$ , such that

$$\begin{array}{ll} \underset{f(x)}{\text{minimize}} & \max_{x \in [a, b]} |f(x)| \\ \text{subject to} & f(0) = 1 \\ & f(x) \text{ is a } T\text{th degree polynomial} \end{array}$$

If  $T = 1$ , then the best function is  $f(x) = 1 - \frac{2}{a+b}x$ .

For higher  $T$ ,  $f(x)$  will be a scaled and shifted Chebyshev polynomial, which gives

$$\alpha(k) = \frac{b-a}{2} \cos\left(\frac{2k+1}{2T}\pi\right) + \frac{a+b}{2}, k = 0, \dots, T-1.$$

## 2 Consensus with Noise

We use the following consensus scheme:

$$x(k+1) = (I - \alpha L)x(k),$$

where  $L$  is the Laplacian matrix and  $\alpha > 0$ . Hence,

$$x_i(k+1) = (1 - d_i\alpha)x_i(k) + \alpha \sum_{j \in \mathcal{N}_i} x_j(k), \quad (2)$$

where  $\mathcal{N}_i$  is the set of the neighboring node of  $i$  and  $d_i = |\mathcal{N}_i|$  is the degree of node  $i$ .

Notice that  $x_j(k)$  in (2) is the message received by node  $i$  from node  $j$ . Now consider that instead of receiving  $x_j(k)$ , the node receives  $z_{ij}(k)$ , which is a noisy version of  $x_j(k)$ :

$$z_{ij}(k) = x_j(k) + w_{ij}(k).$$

Hence, (2) becomes:

$$x_i(k+1) = (1 - d_i\alpha)x_i(k) + \alpha \sum_{j \in \mathcal{N}_i} x_j(k) + \alpha \sum_{j \in \mathcal{N}_i} w_{ij}(k).$$

Define

$$v(k) = \begin{bmatrix} \sum_{j \in \mathcal{N}_1} w_{1j}(k) \\ \vdots \\ \sum_{j \in \mathcal{N}_n} w_{nj}(k) \end{bmatrix}.$$

Therefore,

$$x(k+1) = (I - \alpha L)x(k) + \alpha w(k),$$

where we assume that  $w(k)$  is i.i.d., zero mean and has a bounded second moment. The covariance of  $w(k)$  is defined as  $Q$

If  $\alpha$  is fixed, then we cannot achieve consensus. Hence, we need to use a time varying  $\alpha(k)$ .

$$x(k+1) = (I - \alpha(k)L)x(k) + \alpha(k)v(k), \quad (3)$$

We choose  $\alpha(k) \geq 2/(\lambda_2(L) + \lambda_n(L))$  to satisfies the following condition:

1.  $\sum_{k=0}^{\infty} \alpha(k) = \infty$ .
2.  $\sum_{k=0}^{\infty} \alpha(k)^2 < \infty$ .

Condition 2 implies that  $\alpha(k) \rightarrow 0$ .  
 One possible choice  $\alpha(k) = 1/(k+1)$ .

$$\frac{1}{1^2} + \frac{1}{2^2} + \frac{1}{3^2} = \frac{\pi^2}{6} < 0.$$

In fact, we can choose  $\alpha(k) = (k+1)^{-\varphi}$ , for any  $0.5 < \varphi < 1$ .

Define  $y(k) = x(k) - Jx(k)$ . (Notice that this definition is different from our previous one, where  $y(k) = x(k) - Jx(0)$ . *why?*) Define  $\theta(k) = \mathbf{1}^T x(k)/n$ . Hence,  $x(k) = \theta(k)\mathbf{1} + y(k)$ .

By (3), we have

$$\theta(k+1) = \theta(k) + \alpha(k)\mathbf{1}^T v(k)/n.$$

Hence, for any  $k_1 > k_2$ ,

$$\mathbb{E}(\theta(k_1) - \theta(k_2))^2 = \frac{\mathbf{1}^T Q \mathbf{1}}{n^2} \sum_{t=k_2}^{k_1-1} \alpha(t)^2.$$

Hence,  $\theta(k)$  converges in  $L_2$ . Define  $\theta$  as the  $L_2$  limit of  $\theta_k$ .

Now let us look at  $y(k)$ . By (3),

$$y(k+1) = [(I - \alpha(k)L)(I - J)]y(k) + \alpha(k)(I - J)v(k).$$

Let us define  $\mathcal{P}(k) = (I - \alpha(k)L)(I - J)$ . Therefore,

$$y(k+1) = \prod_{t=0}^k \mathcal{P}(t)y(0) + \sum_{\tau=0}^k \left( \prod_{t=\tau+1}^k \mathcal{P}(t) \right) \alpha(\tau)(I - J)v(\tau).$$

Clearly,

$$\|\mathcal{P}(k)\| = 1 - \alpha(k)\lambda_2(L).$$

Hence,

$$\left\| \prod_{t=k_1}^{k_2} \mathcal{P}(t) \right\| \leq \prod_{t=k_1}^{k_2} (1 - \alpha(t)\lambda_2(L)) \leq \prod_{t=k_1}^{k_2} \exp(-\alpha(t)\lambda_2(L)) = \exp\left(-\sum_{t=k_1}^{k_2} \alpha(t)\lambda_2(L)\right),$$

which implies that

$$\lim_{k \rightarrow \infty} \prod_{t=0}^k \mathcal{P}(k) = 0.$$

On the other hand,

$$\mathbb{E} \left\| \left( \prod_{t=\tau+1}^k \mathcal{P}(t) \right) \alpha(\tau)(I - J)v(\tau) \right\|^2 \leq \beta \exp\left(-2 \sum_{t=\tau+1}^k \alpha(t)\lambda_2(L)\right) \alpha(\tau)^2.$$

where  $\beta = \text{tr}((I - J)Q(I - J))$ . Hence

$$\begin{aligned} \mathbb{E} \left\| \sum_{\tau=0}^k \left( \prod_{t=\tau+1}^k \mathcal{P}(t) \right) \alpha(\tau)(I - J)v(\tau) \right\|^2 &= \sum_{\tau=0}^k \mathbb{E} \left\| \left( \prod_{t=\tau+1}^k \mathcal{P}(t) \right) \alpha(\tau)(I - J)v(\tau) \right\|^2 \\ &\leq \beta \sum_{\tau=0}^k \left[ \exp \left( -2 \sum_{t=\tau+1}^k \alpha(t)\lambda_2(L) \right) \alpha(\tau)^2 \right] \rightarrow 0. \end{aligned}$$

Hence,  $y(k) \rightarrow 0$ . As a result,  $x(k)$  converges to  $\theta \mathbf{1}$  in the mean square sense ( $L_2$ ).