

Thm. 6.4 FloodSet solves the agreement problem in $f+1$ rounds.

PF. Termination ✓

② Validity ✓

③ Agreement

Assume at r' there are no failures
then active P_i 's exchange
and $W_i(r') = W_j(r')$ $\forall i, j$ active.

r remain identical throughout

$\Rightarrow W_i(r+1) = W_j(r+1)$ $\forall i, j$ active
 \therefore decide the same //

Thm 6.33

Cannot guarantee solution of agreement
w/ stopping failures in fewer than $f+1$ rounds

pf method
assume f -round algorithm exists;
and obtain a contradiction
use "